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EU - NCG 4th Annual Meeting in Bucharest, Romania

Background	The Main Idea	The Planar Algebra \mathcal{P}'	The Algorithm	Theorem and Partial Proof	Conclusion
Outline					



- 2 The Main Idea
- **3** The Planar Algebra \mathcal{P}'
- 4 The Algorithm
- 5 Theorem and Partial Proof

6 Conclusion



• Alexander (1923/28) defined Alexander polynomial, *P*(*q*) "Topological Invariants of Knots and Links"



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- Murakami (1993) List of axioms for the MVAP
- Bigelow (Spring 2010) diagrammatic algorithm for the single variable Alexander polynomial using planar algebras



Take knots and links to be oriented tangles with 1 unclosed strand with endpoints on the boundary of a disk

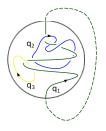
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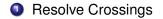




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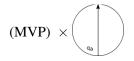
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- Apply Relations

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MVAP=(MVP)(a normalizing coefficient)



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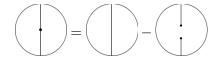


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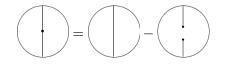
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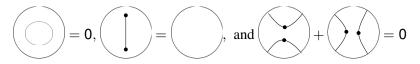


Background The Main Idea The Planar Algebra \mathcal{P}' The Algorithm Theorem and Partial Proof Conclusion Let \mathcal{P}' be the planar algebra generated by a single element in P_1 :

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. The planar algebra \mathcal{P}' has relations in \textit{P}_0 and \textit{P}_4 respectively





Replace the crossings in the link with a linear combination of diagrams in P_4 .



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And you get

$$\Sigma_{k=1}^{5^n} P_k(colors) D_k$$

where n = crossings and each D_k is a diagram in P_2

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Apply relations in the planar algebra \mathcal{P}^\prime and the definition of the dotted strand.

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This will give a multivariable polynomial times a single strand, and this gives the MVAP up to a normalizing coefficient for R1 and the Murakami relations.



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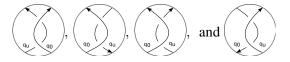
and the Murakami Relations



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$$(\mathbf{x}_{u}) = q_{o}(\mathbf{x}) + q_{o}(\mathbf{x}) + (q_{u} - q_{u}^{-1})(\mathbf{x}) + q_{o}^{-1}(\mathbf{x}) - q_{o}^{-1}(\mathbf{x})$$

$$(\mathbf{x}_{o}) = q_{o} (\mathbf{x}_{o}) + q_{o} (\mathbf{x}_{o}) + (q_{u} - q_{u}^{-1}) (\mathbf{x}_{o}) + q_{o}^{-1} (\mathbf{x}_{o}) - q_{o}^{-1} (\mathbf{x}_{o})$$

$$= q_o \left[q_o^{-1} \left(\begin{array}{c} \\ \end{array} \right) + q_o^{-1} \left(\begin{array}{c} \\ \end{array} \right) - \left(q_u - q_u^{-1} \right) \left(\begin{array}{c} \\ \end{array} \right) + q_o \left(\begin{array}{c} \\ \end{array} \right) - q_o \left(\begin{array}{c} \\ \end{array} \right) \right]$$

$$(\mathbf{x}_{u} = q_{o} (\mathbf{x}_{u} + q_{o} (\mathbf{x}_{u} + q_{u}^{-1}) (\mathbf{x}_{u} + q_{o}^{-1} (\mathbf{x}_{u} - q_{o}^{-1}) (\mathbf{x}_{u} + q_{o}^{-1} (\mathbf{x}_{u} + q_{o}^{-1}))))$$

$$= q_o \left[q_o^{-1} \left(\begin{array}{c} \\ \end{array} \right) + q_o^{-1} \left(\begin{array}{c} \\ \end{array} \right) - \left(q_u - q_u^{-1} \right) \left(\begin{array}{c} \\ \end{array} \right) + q_o \left(\begin{array}{c} \\ \end{array} \right) - q_o \left(\begin{array}{c} \\ \end{array} \right) \right]$$

$$+ q_o \left[q_o^{-1} \cdot 0 + q_o^{-1} \left(\begin{array}{c} \\ \end{array} \right) - \left(q_u - q_u^{-1} \right) \left(\begin{array}{c} \\ \end{array} \right) + \text{ZEROS} \right]$$

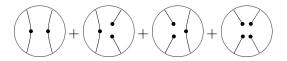
$$(\mathbf{x}_{u}) = q_{o}(\mathbf{x}) + q_{o}(\mathbf{x}) + (q_{u} - q_{u}^{-1})(\mathbf{x}) + q_{o}^{-1}(\mathbf{x}) - q_{o}^{-1}(\mathbf{x})$$

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$$+ (q_u - q_u^{-1}) \Big[q_o (\mathbf{x}) + \text{ZEROS} \Big] + q_0^{-1} \Big[q_o (\mathbf{x}) + \text{ZEROS} \Big] \\ - q_o^{-1} \Big[- q_o (\mathbf{x}) + \text{ZEROS} \Big]$$

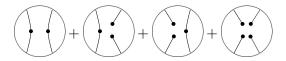
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Conclusion

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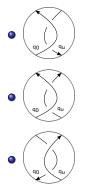


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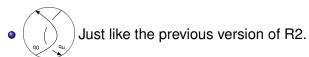
Two iterations of $\bigcirc = \bigcirc - \bigcirc 1$ will finish the verification.

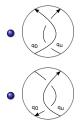
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Other versions of R2

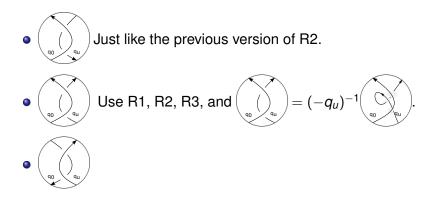




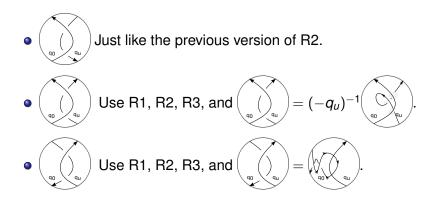












Murakami's Axioms

1

2

5

$$() - () = (q - q^{-1})$$

- Gets its own slide.
- If *L* is the trivial knot with color q_a , then $\Delta(L) = \frac{1}{q_a q_a^{-1}}$.

$$(\bigcirc q_a - q_a^{-1}) \bigcirc q_a$$

• If *L* is the split union of a link and trivial knot, then $\Delta(L)$ is zero.

Murakami's Third Axiom

- Defined in the braid group with three strands, B_3 .
- From left to right the colors of the strands in *B*₃ are *q*₁, *q*₂, and *q*₃.

Murakami's Third Axiom

- Defined in the braid group with three strands, B_3 .
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Define $f_+(x) = x + x^{-1}$ and $f_-(x) = x - x^{-1}$.



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Define
$$f_+(x) = x + x^{-1}$$
 and $f_-(x) = x - x^{-1}$.

$$\begin{split} f_{+}(q_{1})f_{-}(q_{2})\Delta([\sigma_{1}^{-1}\sigma_{2}^{-1}\sigma_{2}^{-1}\sigma_{1}^{-1}]) &- f_{-}(q_{2})f_{+}(q_{3})\Delta([\sigma_{2}^{-1}\sigma_{1}^{-1}\sigma_{1}^{-1}\sigma_{2}^{-1}]) - \\ f_{-}(q_{1}^{-1}q_{3})[\Delta([\sigma_{1}^{-1}\sigma_{1}^{-1}\sigma_{2}^{-1}\sigma_{1}^{-1}]) + \Delta([\sigma_{2}^{-1}\sigma_{2}^{-1}\sigma_{1}^{-1}\sigma_{1}^{-1}])] + f_{-}(q_{1}^{-1}q_{2}q_{3})\Delta([\sigma_{2}^{-1}\sigma_{2}^{-1}]) \\ &- f_{+}(q_{1})f_{-}(q_{1}q_{2}q_{3}^{-1})\Delta([\sigma_{1}^{-1}\sigma_{1}^{-1}]) + f_{-}(q_{1}^{-2}q_{3}^{2})\Delta([e]) \\ &= \text{ZERO} \end{split}$$

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This invariant generalizes easily to tangle invariants up to R1.

- What does the image in \mathcal{P}' of the tangles look like?
- Can this be done for the HOMFLY polynomial?
- Is it possible to find new knot invariants with this method?